

The law of large numbers in categorical probability

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based on work with Tomáš Gonda, Paolo Perrone and Eigil Rischel

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References

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- ▷ Tobias Fritz,
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Adv. Math. 370, 107239 (2020). [arXiv:1908.07021](#).
- ▷ Tobias Fritz and Eigil Fjeldgren Rischel,
The zero-one laws of Kolmogorov and Hewitt–Savage in categorical probability.
Compositionality 2, 3 (2020). [arXiv:1912.02769](#).
- ▷ Tobias Fritz, Tomáš Gonda, Paolo Perrone,
De Finetti’s Theorem in Categorical Probability.
[arXiv:2105.02639](#).
- ▷ ...?

For a broader perspective, see the videos from the online workshop [Categorical Probability and Statistics!](#)

The law of large numbers

Theorem

Let $(x_n)_{n \in \mathbb{N}}$ be a sequence of real-valued independent random variables with identical distribution and $\mathbb{E}[|x_1|] < \infty$. Then

$$\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n x_i = \mathbb{E}[x_1]$$

with probability 1.

- ▷ **Example:** Upon repeatedly tossing a fair coin, the relative frequency of heads approaches $\frac{1}{2}$ with probability 1.
- ▷ But where are the categories?

Teaser

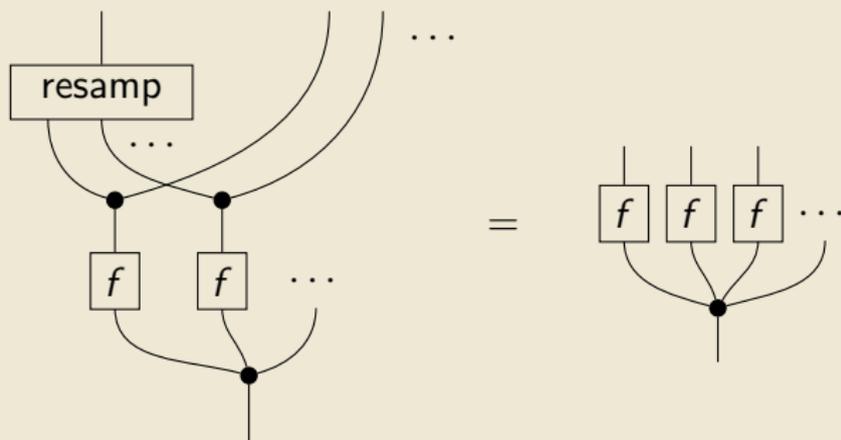
My goal is to explain this form of the law of large numbers:

Theorem/Definition

For every object X there is a partial morphism

$$\text{resamp} : X^{\mathbb{N}} \rightarrow X$$

such that for every $f : A \rightarrow X$,



The big picture

Traditional probability theory	Categorical probability theory
Analytic: says what probabilities are	Synthetic: says how probabilities behave
Analogous to number systems	Analogous to abstract algebra

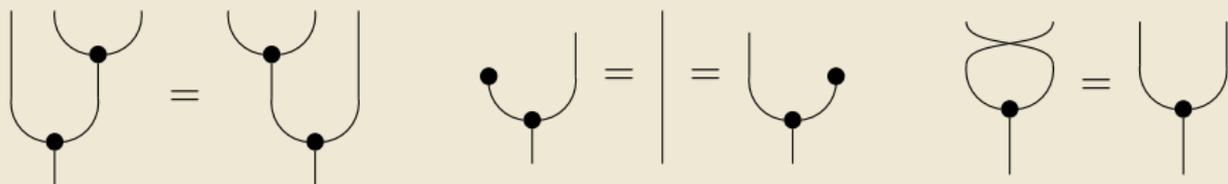
- ▷ There will be no numerical probabilities!
- ▷ I can say more about the motivations and scope of categorical probability if there is need for discussion.

Definition

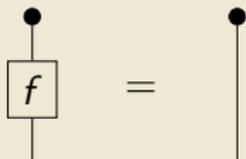
A **Markov category** is a symmetric monoidal category supplied with **copying** and **deleting** operations on every object,



giving commutative comonoid structures



which interact well with the monoidal structure, and such that for all f ,



Semantics

There are many different (and interesting) Markov categories.

But for today, I have one particular intended semantics in mind:

Definition

BorelStoch is the category with:

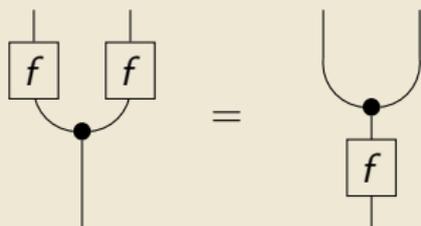
- ▶ **Standard Borel spaces** as objects (finite sets, \mathbb{N} and $[0, 1]$).
- ▶ Measurable Markov kernels as morphisms.
- ▶ Products of measurable spaces for \otimes .

BorelStoch encodes standard measure-theoretic probability.

Determinism

Definition

A morphism $f : X \rightarrow Y$ is **deterministic** if it commutes with copying,



- ▷ **Intuition:** Applying f to copies of input = copying the output of f .
- ▷ Deterministic morphisms form a cartesian monoidal subcategory \mathbf{C}_{det} .
- ▷ $\mathbf{BorelStoch}_{\text{det}}$ is the category of measurable functions between standard Borel spaces.

Infinite tensor products

Let $(X_n)_{n \in \mathbb{N}}$ be a family of objects.

For finite $F \subseteq F' \subseteq \mathbb{N}$, we have projection morphisms

$$\bigotimes_{n \in F'} X_n \longrightarrow \bigotimes_{n \in F} X_n$$

given by composing with deletion for all $n \in F' \setminus F$, like this:

$$\begin{array}{ccc} X_1 & & \\ | & \cdots & \bullet \\ X_1 & & X_n \end{array}$$

Infinite tensor products

Definition

The **infinite tensor product**

$$X^{\mathbb{N}} = \bigotimes_{n \in \mathbb{N}} X_n$$

is the limit of the finite tensor products $X^F := \bigotimes_{n \in F} X_n$ if it exists and is preserved by every $- \otimes Y$.

- ▷ **Intuition:** To map into an infinite tensor product, one needs to map consistently into its finite subproducts.

Kolmogorov products

Definition

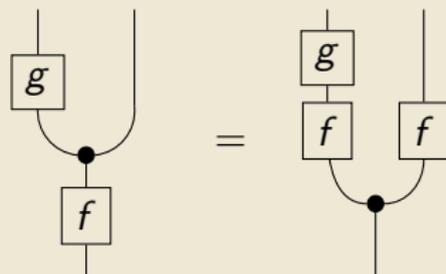
An infinite tensor product $X^{\mathbb{N}}$ is a **Kolmogorov product** if the limit projections $\pi^F : X^{\mathbb{N}} \rightarrow X^F$ are deterministic.

- ▷ This additional condition fixes the comonoid structure on $X^{\mathbb{N}}$.
- ▷ From now on: assume Markov category with countable Kolmogorov products.
- ▷ Satisfied by **BorelStoch** (Kolmogorov extension theorem).

Positivity

Definition

\mathbf{C} is **positive** if the following holds: if a composite gf is deterministic, then also



- ▷ **Intuition:** If a deterministic process has a random intermediate result, then that result can be computed independently from the process.
- ▷ Positivity implies that every isomorphism is deterministic.
- ▷ Not every Markov category is positive.

Partial morphisms

- ▷ In the law of large numbers, the limit

$$\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n x_i$$

does not always exist.

- ▷ This suggests the need for **partial morphisms** in categorical probability.
- ▷ Under certain “partializability” conditions we indeed get a **monoidal restriction category**.

Partial morphisms

Definition

A positive Markov category is **partializable** if deterministic monos are closed under

- ▷ pullbacks,
 - ▷ tensor products.
-
- ▷ In **BorelStoch**, the deterministic subobjects of X are the **measurable sets** $S \subseteq X$.
 - ▷ This is a deep fact of descriptive set theory!
 - ▷ Thanks to it, one can show that **BorelStoch** is partializable.

Resampling

- ▷ In statistics, **resampling** is a set of methods to estimate generalization (e.g. cross-validation).
- ▷ Here, I mean something different but closely related:

Resampling = Pick an element from an infinite sequence uniformly at random
--

- ▷ This doesn't make literal sense since **there is no uniform distribution on \mathbb{N}** .
- ▷ But it can be made to work for *some* sequences (x_n) .
- ▷ Thus we get a partial morphism

$$\text{resamp} : X^{\mathbb{N}} \rightarrow X.$$

Resampling

- ▷ In **BorelStoch**, we would want to define

$$\text{resamp}(S \mid (x_n)) := \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n 1_S(x_i)$$

whenever this limit exists for all measurable S .

- ▷ For finite n , this indeed corresponds to choosing an element from a finite sequence uniformly at random.
- ▷ The problem is that the resulting

$$\text{resamp}(- \mid (x_n))$$

would not always be a probability measure.

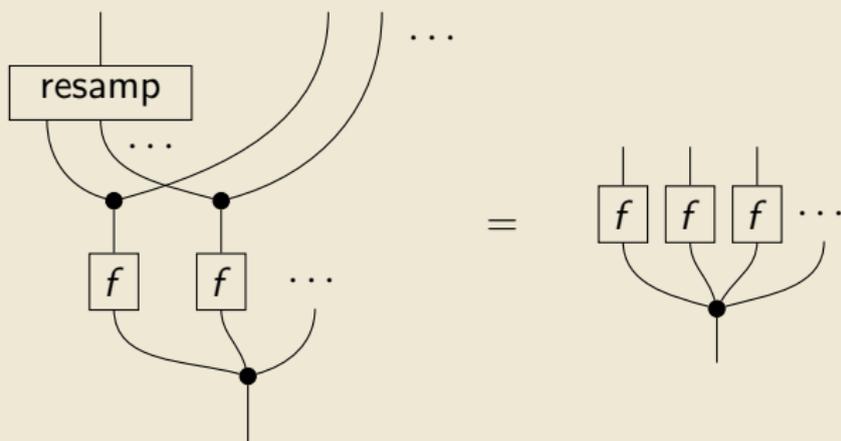
- ▷ Can be fixed by imposing **uniform** existence of the limits.

Back to the law of large numbers

We can now understand the following:

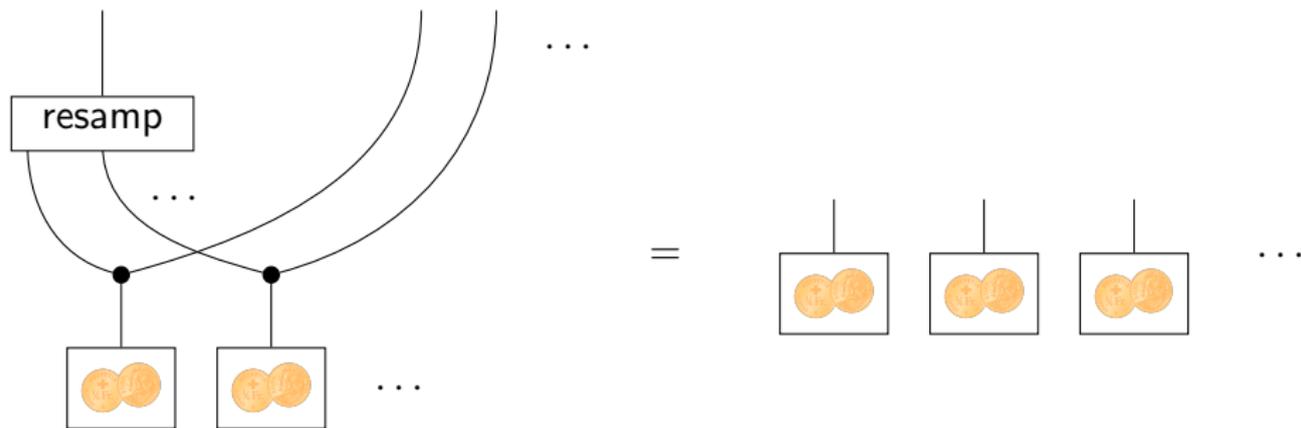
Theorem

In **BorelStoch**, every $f : A \rightarrow X$ satisfies



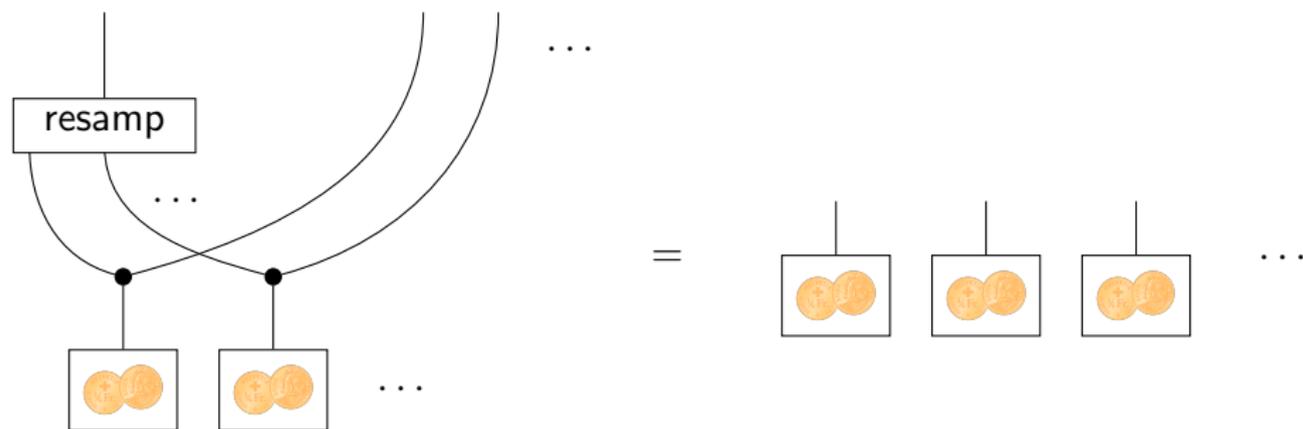
This statement encodes the **Glivenko–Cantelli theorem** on the convergence of the empirical distribution, a strong law of large numbers.

Coin example



- ▷ **Interpretation:** flipping infinitely many fair coins and then picking a random one makes the latter
 - ▷ independent of, and
 - ▷ identically distributed asthe others.

Coin example



▷ In terms of random outcomes $c_i \in \{\text{heads}, \text{tails}\}$, this equation says

$$\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n 1_{\text{heads}}(c_i) =_{\text{a.s.}} \mathbb{P}[\text{heads}] = \frac{1}{2},$$

an instance of the law of large numbers.

Summary

- ▷ Markov categories = emerging framework for synthetic probability.
- ▷ We have abstract versions of some theorems of probability and statistics:
 - ▷ 0/1-laws of Kolmogorov and Hewitt-Savage,
 - ▷ Fisher factorization theorem on sufficient statistics,
 - ▷ Blackwell-Sherman-Stein theorem on informativeness of statistical experiments,
 - ▷ de Finetti theorem on exchangeable distributions.
- ▷ We should try to add the law of large numbers to this list!
- ▷ There are further hints of connections with ergodic theory.

Bonus slides: Why categorical probability?

In no particular order:

- ▷ Applications to probabilistic programming.
- ▷ Prove theorems in greater generality and with more intuitive proofs.
- ▷ Reverse mathematics: sort out interdependencies between theorems.
- ▷ Ultimately, prove theorems of higher complexity?
- ▷ Simpler teaching of probability theory. (String diagrams!)
- ▷ Different conceptual perspective on what probability is.

Discrete probability theory as a Markov category

One of the paradigmatic Markov categories is **FinStoch**, the category of finite sets and **stochastic matrices**: a morphism $f : X \rightarrow Y$ is

$$(f(y|x))_{x \in X, y \in Y} \in \mathbb{R}^{X \times Y}$$

with

$$f(y|x) \geq 0, \quad \sum_y f(y|x) = 1.$$

Composition is the **Chapman-Kolmogorov formula**,

$$(gf)(z|x) := \sum_y g(z|y) f(y|x).$$

A morphism $p : 1 \rightarrow X$ is a **probability distribution**.

A general morphism $X \rightarrow Y$ has many names: **Markov kernel**, **probabilistic mapping**, **communication channel**, ...

The monoidal structure implements **stochastic independence**,

$$(g \otimes f)(xy|ab) := g(x|a) f(y|b).$$

The copy maps are

$$\text{copy}_X : X \longrightarrow X \times X, \quad \text{copy}_X(x_1, x_2|x) = \begin{cases} 1 & \text{if } x_1 = x_2 = x, \\ 0 & \text{otherwise.} \end{cases}$$

The deletion maps are the unique morphisms $X \rightarrow 1$.

- ▷ Works just the same with “probabilities” taking values in any **semiring** R .
- ▷ Taking R to be the **Boolean semiring** $\mathbb{B} = \{0, 1\}$ with

$$1 + 1 = 1$$

results in the Kleisli category of the nonempty finite powerset monad.

⇒ We get a Markov category for non-determinism.

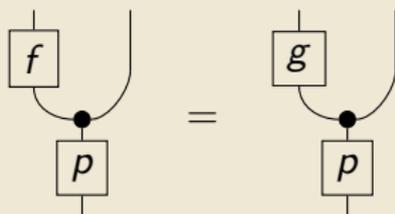
- ▷ Measure-theoretic probability: Kleisli category of the **Giry monad**.

Almost sure equality

Definition

Let $p : A \rightarrow X$ and $f, g : X \rightarrow Y$.

f and g are **equal p -almost surely**, $f =_{p\text{-a.s.}} g$, if

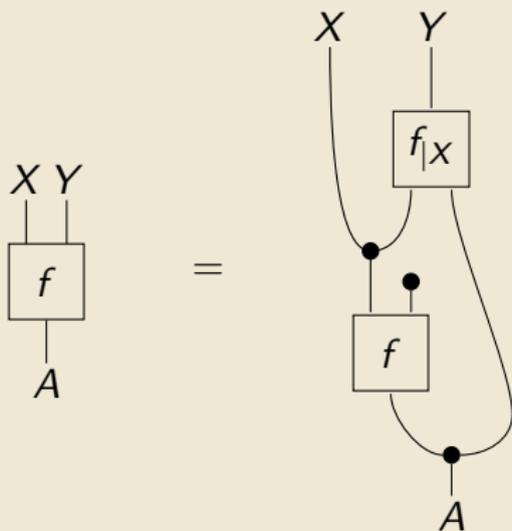


- ▷ **Intuition:** f and g behave the same on all inputs produced by p .
- ▷ In **BorelStoch**, coincides with the standard notion of a.s. equality.
- ▷ Other concepts relativize similarly with respect to p -almost surely.

Conditionals

Definition

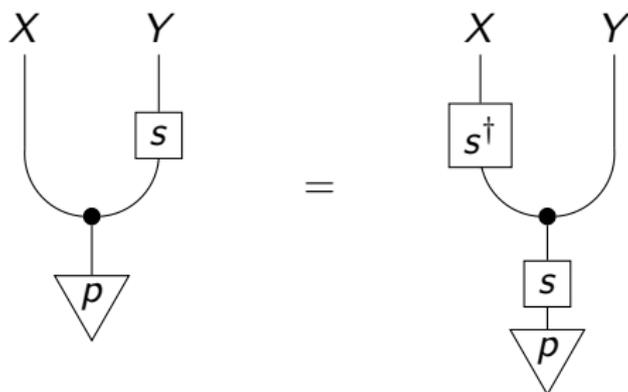
A Markov category has conditionals if for every $f : A \rightarrow X \otimes Y$ there is $f_{|X} : X \otimes A \rightarrow Y$ with



▷ **Intuition:** The outputs of f can be generated one at a time.

Bayesian inversion

Every $s : X \rightarrow Y$ has a **Bayesian adjoint** $s^\dagger : Y \rightarrow X$ satisfying:

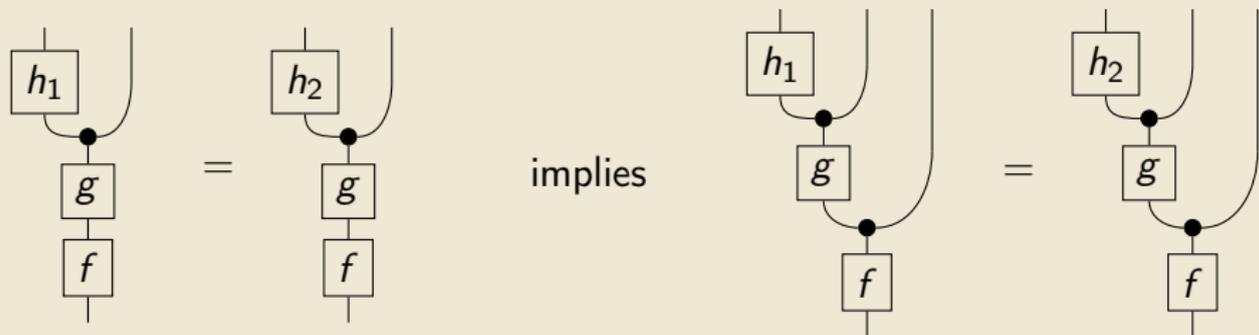


The Bayesian adjoint s^\dagger depends on p .

The causality axiom

Definition

C is causal if

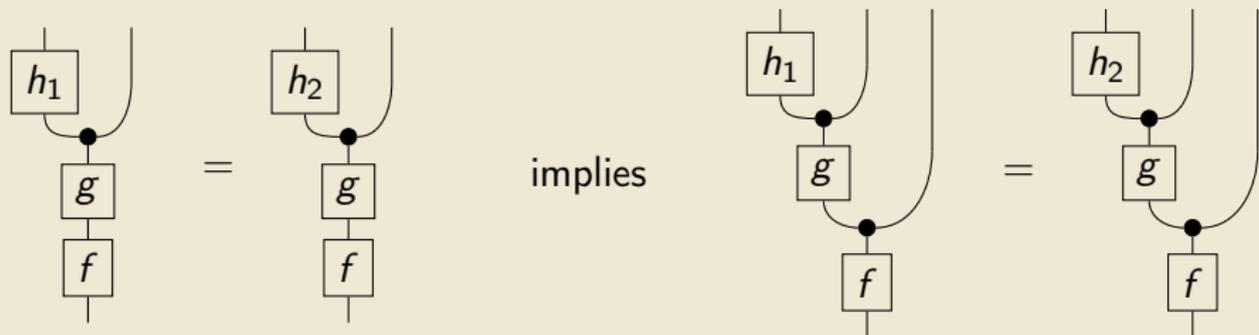


- ▷ **Intuition:** The choice between h_1 and h_2 in the “future” of g does not influence the “past” of g .
- ▷ Not every Markov category is causal.

The causality axiom

Definition

C is causal if



- ▷ **Intuition:** The choice between h_1 and h_2 in the “future” of g does not influence the “past” of g .
- ▷ Not every Markov category is causal.

Representability

Definition

A Markov category \mathbf{C} is **representable** if for every $X \in \mathbf{C}$ there is $PX \in \mathbf{C}$ and a natural bijection

$$\mathbf{C}_{\text{det}}(-, PX) \cong \mathbf{C}(-, X),$$

and **a.s.-compatibly representable** if this respects p -a.s. equality for every p .

- ▶ **Intuition:** PX is space of probability measures on X .
- ▶ Under the bijection, the deterministic $\text{id} : PX \rightarrow PX$ corresponds to

$$\text{samp}_X : PX \rightarrow X,$$

the map that returns a random sample from a distribution.

Kleisli categories are Markov categories

Proposition

Let

- ▷ \mathbf{D} be a category with finite products,
- ▷ P a commutative monad on \mathbf{D} with $P(1) \cong 1$.

Then the Kleisli category $\text{Kl}(P)$ is a Markov category in the obvious way.

Examples:

- ▷ Kleisli category of the Giry monad, other related monads for measure-theoretic probability.
- ▷ Kleisli category of the non-empty power set monad, which is (almost) **Rel**.

The proposition still holds when \mathbf{D} is merely a Markov category itself!

Categories of comonoids

Proposition

Let \mathbf{C} be any symmetric monoidal category. Then the category with:

- ▷ Commutative comonoids in \mathbf{C} as objects,
- ▷ Counital maps as morphisms,
- ▷ The specified comultiplications as copy maps,

is a Markov category.

A good example is $\mathbf{Vect}_k^{\text{op}}$ for a field k :

- ▷ The comonoids correspond to commutative k -algebras of k -valued random variables.
- ▷ We obtain **algebraic probability theory** with “random variable transformers” as morphisms (formal opposites of Markov kernels).

Diagram categories and ergodic theory

Proposition

Let \mathbf{D} be any category and \mathbf{C} a Markov category. The category in which

- ▷ Objects are functors $\mathbf{D} \rightarrow \mathbf{C}_{\text{det}}$,
- ▷ Morphisms are natural transformations with components in \mathbf{C} .

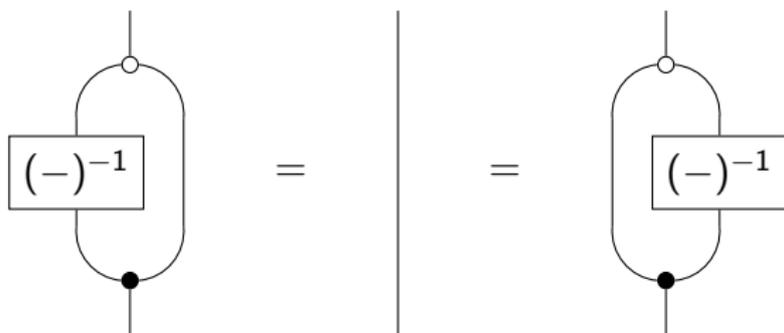
With the poset $\mathbf{D} = \mathbb{Z}$, we get a category of **discrete-time stochastic processes**.

This generalizes an observation going back to (Lawvere, 1962).

We can also take $\mathbf{D} = \mathbf{B}G$ for a group G , resulting in categories of dynamical systems with deterministic dynamics but stochastic morphisms.

Hyperstructures: categorical algebra in Markov categories

A **group** G is a monoid G together with $(-)^{-1} : G \rightarrow G$ such that



This equation can be interpreted in any Markov category! (Together with the bialgebra law.)

- ▷ More generally, one can consider models of any algebraic theory in any Markov category.
- ▷ In Kleisli categories of probability-like monads, these are known as **hyperstructures**.
- ▷ Peter Arndt's suggestion:
Develop categorical algebra for hyperstructures in terms of Markov categories!

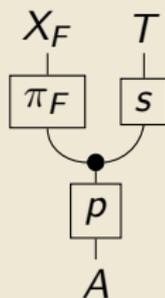
The synthetic Kolmogorov zero–one law

Theorem

Let X_I be a Kolmogorov product of a family $(X_i)_{i \in I}$.

If

- ▷ $p : A \rightarrow X_I$ makes the X_i independent and identically distributed, and
- ▷ $s : X_I \rightarrow T$ is such that



displays $X_F \perp T \parallel A$ for every finite $F \subseteq I$,

then ps is deterministic.

The classical Hewitt–Savage zero-one law

Theorem

Let $(x_n)_{n \in \mathbb{N}}$ be independent and identically distributed random variables, and S any event depending only on the x_n and invariant under finite permutations.

Then $P(S) \in \{0, 1\}$.

The synthetic Hewitt–Savage zero-one law

Theorem

Let J be an infinite set and \mathbf{C} a causal Markov category. Suppose that:

- ▷ The Kolmogorov power $X^{\otimes J} := \lim_{F \subseteq J \text{ finite}} X^{\otimes F}$ exists.
- ▷ $p : A \rightarrow X^{\otimes J}$ displays the conditional independence $\perp_{i \in J} X_i \parallel A$.
- ▷ $s : X^J \rightarrow T$ is deterministic.
- ▷ For every finite permutation $\sigma : J \rightarrow J$, permuting the factors $\tilde{\sigma} : X^{\otimes J} \rightarrow X^{\otimes J}$ satisfies

$$\tilde{\sigma} p = p, \quad s \tilde{\sigma} = s.$$

Then sp is deterministic.

Proof is by string diagrams, but far from trivial!

Detour: random measures

▷ Suppose that I hand you a coin (which may be biased).

▷ How much would you bet on the outcome

heads, tails, tails

when the coin is flipped 3 times?

⇒ Surely the same as you would bet on

tails, tails, heads.

▷ Your bets satisfy **permutation invariance**. Can we say more?

Classical de Finetti theorem

A sequence $(x_n)_{n \in \mathbb{N}}$ of random variables on a space X is **exchangeable** if their distribution is invariant under finite permutations σ ,

$$\begin{aligned} & \mathbb{P} \llbracket x_1 \in S_{\sigma(1)}, \dots, x_n \in S_{\sigma(n)} \rrbracket \\ &= \mathbb{P} \llbracket x_1 \in S_1, \dots, x_n \in S_n \rrbracket. \end{aligned}$$

Theorem

If (x_n) is exchangeable, then there is a measure μ on PX such that

$$\mathbb{P} \llbracket x_1 \in S_1, \dots, x_n \in S_n \rrbracket = \int p(x_1 \in S_1) \cdots p(x_n \in S_n) \mu(dp).$$

Idea: sequence of tosses of a coin with unknown bias!

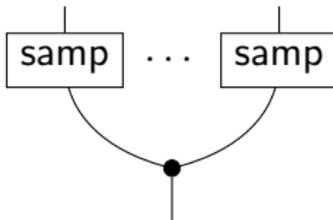
The de Finetti theorem

Assumption: All three axioms above hold. (True for **BorelStoch.**)

Definition

$p : A \rightarrow X^{\mathbb{N}}$ is **exchangeable** if it is invariant under composing with finite permutations.

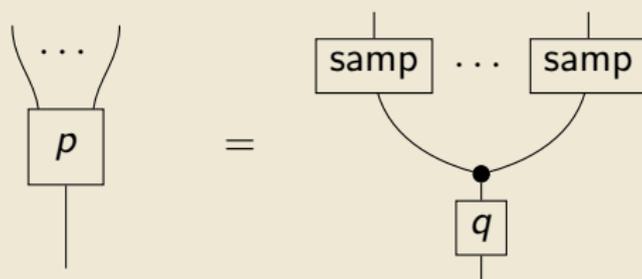
Sampling \mathbb{N} times gives a morphism $PX \rightarrow X^{\mathbb{N}}$ given by



The de Finetti theorem

Theorem

For every exchangeable $p : A \rightarrow X^{\mathbb{N}}$ there is $q : A \rightarrow PX$ such that



- ▷ **Intuition:** The probabilities associated to your bets arise from sampling from a random distribution.